

# Discrete Cosine Transformation & Local Computation

M. Pouly

Department of Informatics  
University of Fribourg, Switzerland

February 2007

# JPEG Image Compression

ISO 10918-1 norm with 41 sub-file-formats. But only one is practically relevant.

Procedure involves 5 steps:

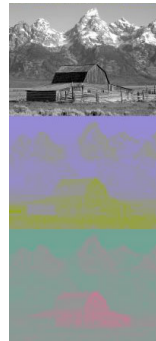
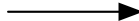
- 1 Color Space Transformation: RGB  $\rightarrow$  YCbCr model
- 2 Downsampling (**lossy**)
- 3 Block Splitting & Discrete Cosine Transformation (DCT)
- 4 Quantization (**lossy**)
- 5 Resorting & Entropy Encoding (Huffman)

top-down = JPEG encoding

bottom-up (with inverse DCT) = JPEG decoding

# Color Space Transformation I

- Image exists in RGB (Red, Green, Blue).
- JPEG uses YCbCr (Y = Pixel brightness, Cb & Cr = Pixel chrominance split into blue and red).



# Color Space Transformation II

This is a linear transformation (CCIR 601 norm):

$$\begin{bmatrix} Y \\ Cb \\ Cr \end{bmatrix} \approx \begin{bmatrix} 0 \\ 128 \\ 128 \end{bmatrix} + \begin{bmatrix} 0.299 & 0.587 & 0.114 \\ -0.168736 & -0.331264 & 0.5 \\ 0.5 & -0.418688 & -0.081312 \end{bmatrix} \cdot \begin{bmatrix} R \\ G \\ B \end{bmatrix}$$

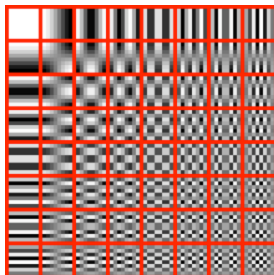
$Y, Cb, Cr \in \{0, \dots, 255\}$ .

# Downsampling

- Humans can see considerably more fine detail in the brightness of an image (Y component) than in the color of an image (Cb and Cr components).
- Therefore, Cr and Cb are stored in a reduced resolution.
- Downsampling the Cb and Cr components saves 33% to 50% of the space taken by the image without drastically affecting perceptual image quality.

## Block Splitting & DCT

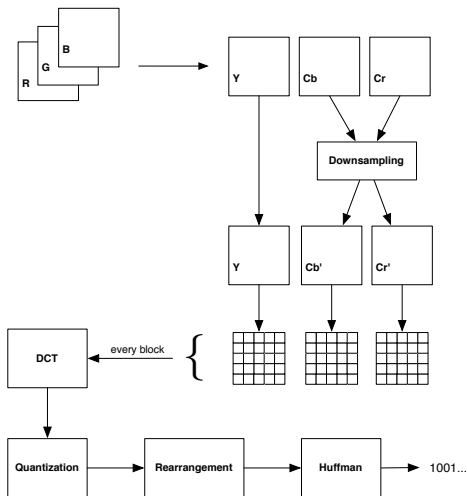
- The Y, Cr and Cb components are split into  $8 \times 8$  blocks.
- We apply a 2-dimensional DCT on each block.
- Instead of 64 points, every block is represented by a linear combination of 64 blocks:



# Quantization & Huffman Encoding

- Human eye is not very sensible to high frequency variations.
- We divide every component in the frequency domain (DCT coefficients) by a constant and round to the nearest integer:
  - many of the higher frequencies become zero
  - others become small positive and negative integers.
- Integers are rearranged and Huffman encoded.

# Summary: The Big Picture



# Discrete Cosine Transformation

$$\mathcal{F}(k) = 2 \sum_{n=0}^{N-1} f(n) \cdot \cos \left[ \frac{\pi k}{N} \cdot \left( n + \frac{1}{2} \right) \right].$$

for  $k = 0, \dots, N - 1$ .

- $f$ : signal ( $\mathbb{C}$ -valued function of discrete time)  $f : \mathbb{Z} \rightarrow \mathbb{C}$
- $N$ : number of signal samples

For Image Compression:

- $f(n) \in \{0, \dots, 255\}$   
is the value of either Y, Cb or Cr at pixel  $n$ .

## Why DCT and not DFT ?

$$\sum_{n=0}^{N-1} f(n) \cdot \cos \left[ \frac{\pi k}{N} \cdot \left( n + \frac{1}{2} \right) \right] \quad \text{versus} \quad \sum_{n=0}^{N-1} f(n) \cdot e^{-\frac{2\pi i n k}{N}}$$

- Transforms images efficiently into a compressible form.
- Uses only **cos**, DFT uses **sin** and **cos** (hidden within  $e^{-i\varphi}$ ).
- DCT coefficients are real, DFT coefficients are complex.

# Towards a Projection Problem I

We use the definition

$$\cos(\varphi) = \frac{1}{2} \left( e^{i\varphi} + e^{-i\varphi} \right)$$

and obtain

$$\begin{aligned} \mathcal{F}(k) &= 2 \sum_{n=0}^{N-1} f(n) \cdot \frac{1}{2} \left( e^{\frac{i\pi k}{N} \cdot (n+\frac{1}{2})} + e^{-\frac{i\pi k}{N} \cdot (n+\frac{1}{2})} \right) \\ &= \sum_{n=0}^{N-1} f(n) \cdot \left( e^{\frac{i\pi kn}{N}} \cdot e^{\frac{i\pi k}{2N}} + e^{-\frac{i\pi kn}{N}} \cdot e^{-\frac{i\pi k}{2N}} \right) \\ &= \sum_{n=0}^{N-1} f(n) \cdot e^{\frac{i\pi kn}{N}} \cdot e^{\frac{i\pi k}{2N}} + \sum_{n=0}^{N-1} f(n) \cdot e^{-\frac{i\pi kn}{N}} \cdot e^{-\frac{i\pi k}{2N}}. \end{aligned}$$

## Towards a Projection Problem II

$$\sum_{n=0}^{N-1} f(n) \cdot e^{\frac{i\pi kn}{N}} \cdot e^{\frac{i\pi k}{2N}} \quad \text{Where are the variables ?}$$

Binary encoding: assume  $N = 2^m$

$$n = \sum_{j=0}^{m-1} n_j \cdot 2^j \quad \text{and} \quad k = \sum_{l=0}^{m-1} k_l \cdot 2^l$$

Thus,  $n = (n_0, \dots, n_{m-1})$  and  $k = (k_0, \dots, k_{m-1})$

$$nk = \sum_{0 \leq j, l \leq m-1} n_j k_l \cdot 2^{j+l}$$

## Towards a Projection Problem III

$$\begin{aligned}
 & \sum_{n=0}^{N-1} f(n) \cdot e^{\frac{i\pi kn}{N}} \cdot e^{\frac{i\pi k}{2N}} = \\
 & \sum_{n_0, \dots, n_{m-1}} f(n_0, \dots, n_{m-1}) \cdot e^{\frac{i\pi \sum_{0 \leq j, l \leq m-1} n_j k_l \cdot 2^{j+l}}{2^m}} \cdot e^{\frac{i\pi \sum_{0 \leq l \leq m-1} k_l \cdot 2^l}{2^{m+1}}} = \\
 & \sum_{n_0, \dots, n_{m-1}} f(n_0, \dots, n_{m-1}) \prod_{0 \leq j, l \leq m-1} e^{\frac{i\pi n_j k_l \cdot 2^{j+l}}{2^m}} \prod_{0 \leq l \leq m-1} e^{\frac{i\pi k_l \cdot 2^l}{2^{m+1}}} = \\
 & \sum_{n_0, \dots, n_{m-1}} f(n_0, \dots, n_{m-1}) \prod_{0 \leq j, l \leq m-1} e^{\frac{i\pi n_j k_l}{2^{m-j-l}}} \prod_{0 \leq l \leq m-1} e^{\frac{i\pi k_l}{2^{m+1-l}}}.
 \end{aligned}$$

# DCT as Projection Problem

$$\sum_{n_0, \dots, n_{m-1}} f(n_0, \dots, n_{m-1}) \prod_{0 \leq j, l \leq m-1} e^{\frac{i\pi n_j k_l}{2^{m-j-l}}} \prod_{0 \leq l \leq m-1} e^{\frac{i\pi k_l}{2^{m+1-l}}} +$$

$$\sum_{n_0, \dots, n_{m-1}} f(n_0, \dots, n_{m-1}) \prod_{0 \leq j, l \leq m-1} e^{-\frac{i\pi n_j k_l}{2^{m-j-l}}} \prod_{0 \leq l \leq m-1} e^{-\frac{i\pi k_l}{2^{m+1-l}}}$$

This are two projection problems !

- Total domain:  $\{k_0, \dots, k_{m-1}, n_0, \dots, n_{m-1}\}$ .
- All variables are binary.
- Query:  $\{k_0, \dots, k_{m-1}\}$ , factor domains:  $\{n_0, \dots, n_{m-1}\}, \{n_j, k_l\}$ .
- Factors:  $f(n_0, \dots, n_{m-1})$  and  $e^{\pm \frac{i\pi n_j k_l}{2^{m-j-l}}}$  and  $e^{\pm \frac{i\pi k_l}{2^{m+1-l}}}$
- Factors are valuations induced by the Arithmetic Semiring.

# Conclusion

- DCT can be written as two projection problem.
  - both queries are identical,
  - both knowledgebases have identical factor domains.
- Ergo: we can use the same covering join tree.